

CSE 260M / ESE 260

Intro. To Digital Logic & Computer Design

Bill Siever
&
Jim Feher

This week

- Thursday:
Studio — Here / Seigle 301
- Hw#5 posted by late Friday
 - Verilog!

Chapter 4

**Studio / Environment: nand2 / add2
(Simulation)**

always: Based on *Events*

- Concept of “event” is related to simulation and “event driven programming”
- JLS uses events: An OR gate “reacts” to events and schedules an update
See [here](#)

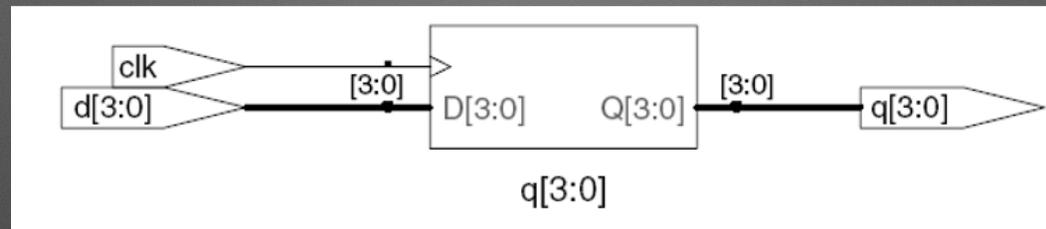
always Statement

- Form:

```
always @(sensitivity list)
    statement;
```

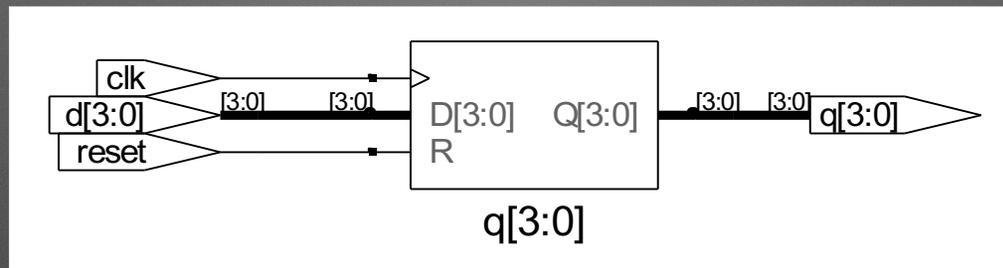
- When event in sensitivity list occurs, statement is executed

Verilog: D Flip-Flop



```
module flop(input logic clk,  
            input logic [3:0] d,  
            output logic [3:0] q);  
    always_ff @(posedge clk)  
        q <= d; // pronounced "q gets d"  
endmodule
```

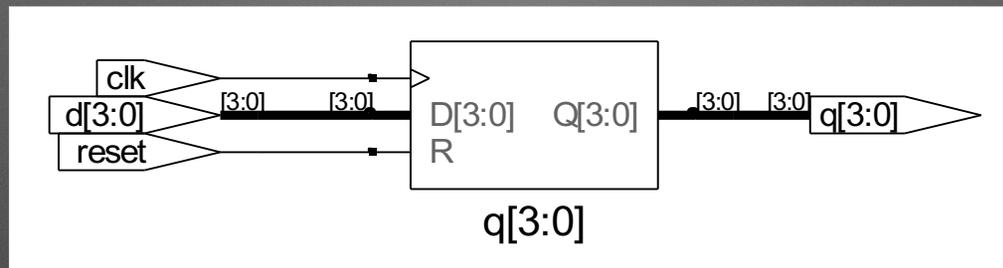
Resettable D-Flip-Flop 1



```
module flopr(input logic clk,
            input logic reset,
            input logic [3:0] d,
            output logic [3:0] q);

    always_ff @(posedge clk)
        if (reset) q <= 4'b0;
        else q <= d;
endmodule
```

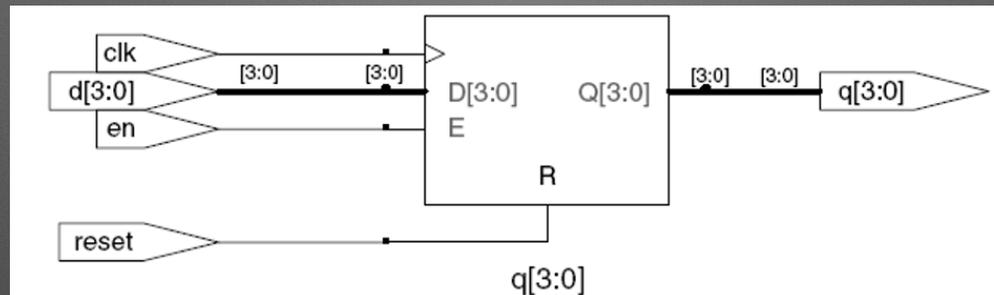
Resettable D-Flip-Flop 2



```
module flopr(input logic clk,
             input logic reset,
             input logic [3:0] d,
             output logic [3:0] q);

    always_ff @(posedge clk, posedge reset)
        if (reset) q <= 4'b0;
        else q <= d;
endmodule
```

Resettable D-Flip-Flop 3



```
module flopr(input logic clk,
            input logic reset,
            input logic en,
            input logic [3:0] d,
            output logic [3:0] q);

    always_ff @(posedge clk, posedge reset)
        if (reset) q <= 4'b0;
        else if (en) q <= d;
endmodule
```

Blocking vs. Non-Blocking Assignment (in always blocks)

- `<=` is nonblocking assignment
Occurs simultaneously with others
- `=` is blocking assignment
Occurs in order it appears in file

Rules for Assignments

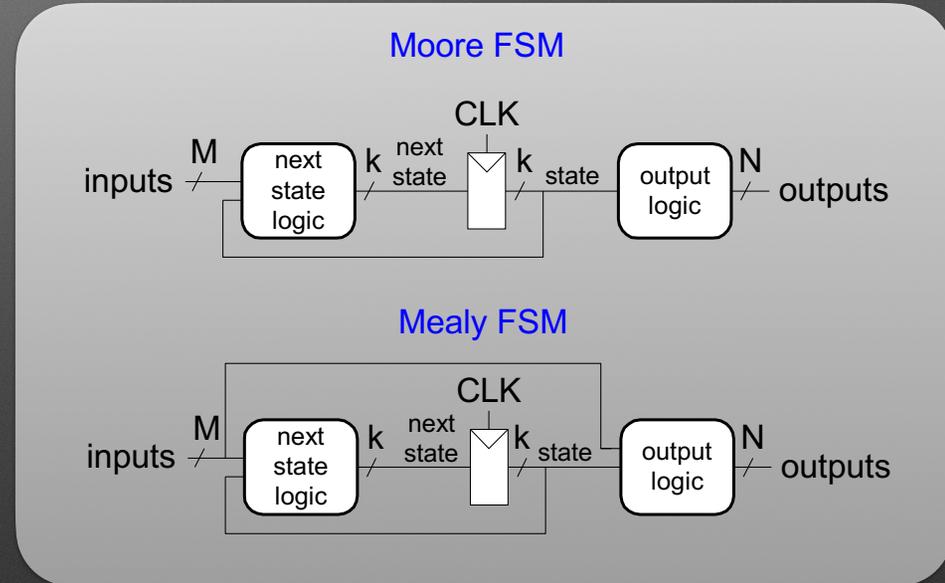
- **Synchronous sequential logic**
use `always_ff @(posedge clk)` and **nonblocking assignments** (`<=`)

```
always_ff @(posedge clk)
    q <= d; // nonblocking
```
- **Simple combinational logic**
use **continuous assignments** (`assign`)

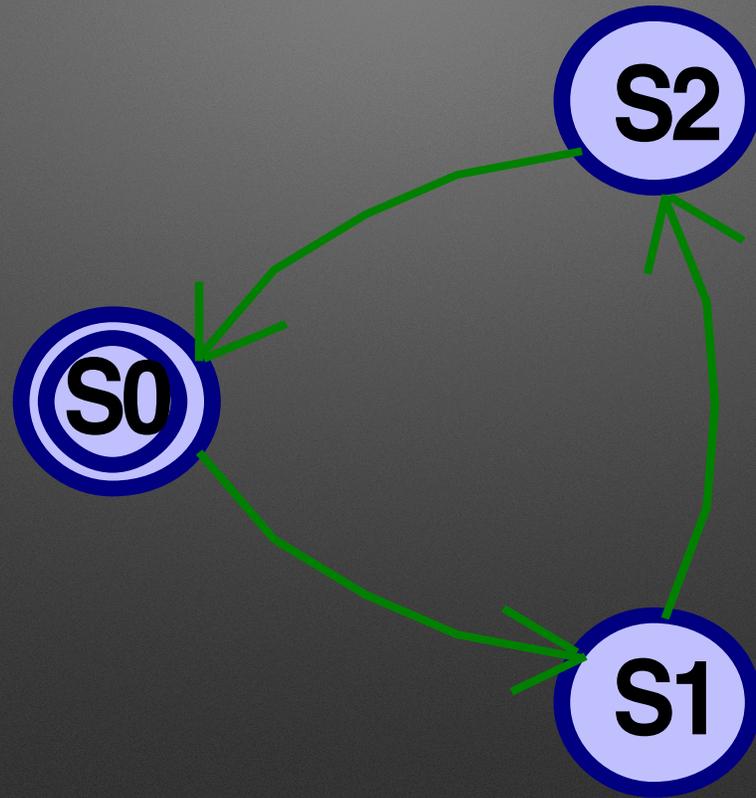
```
assign y = a & b;
```
- **Complex Combinational Logic**
use `always_comb` and **blocking assignments** (`=`)
- **Assign signals in only one `always` or `assign` statement!**

Verilog FSMs

- Three parts
 - Next state logic (arrows / next state table)
 - State register (active bubble)
 - Output logic (output equations)



Divide by 3 Counter



Verilog

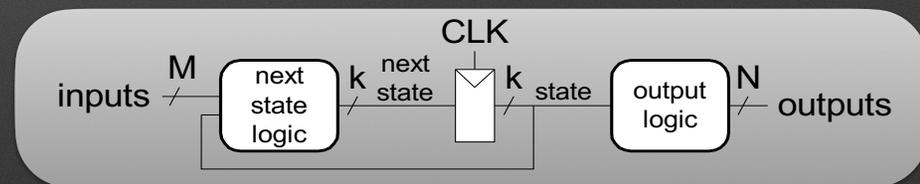
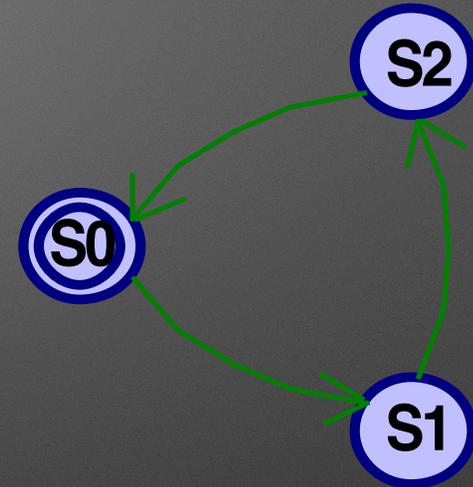
```
module divideby3FSM(input  logic clk,
                   input  logic reset,
                   output logic q);

typedef enum logic [1:0] {S0, S1, S2} statetype;
statetype state, nextstate;

// state register
always_ff @(posedge clk, posedge reset)
    if (~reset) state <= S0;
    else       state <= nextstate;

// next state logic
always_comb
    case (state)
        S0: nextstate = S1;
        S1: nextstate = S2;
        S2: nextstate = S0;
        default: nextstate = S0;
    endcase

// output logic
assign q = (state == S0);
endmodule
```



Parameterized Modules: Declaration

- Way to specify additional details for an *instance* of a generic part
 - Commonly the “width” of the part

```
module mux2
  #(parameter width = 8) // name and default value
  (input  logic [width-1:0] d0, d1,
   input  logic           s,
   output logic [width-1:0] y);
  assign y = s ? d1 : d0;
endmodule
```

Parameterized Modules: Use

- Default or specify parameter for instance:

```
mux2 myMux(d0, d1, s, out);
```

```
mux2 #(12) lowmux(d0, d1, s, out);
```

Ports: Positional vs. Named

- Default or specify parameter for instance:

```
logic a, b, sel, y
mux2 myMux(a, b, sel, y);
```

vs.

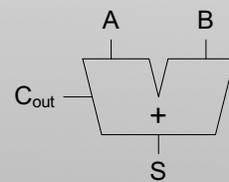
```
mux2 myMux(.d0(a), .d1(b),
           .s(sel), .out(y));
```

```
module mux2
    #(parameter width = 8)
    (input logic [width-1:0] d0,
     input logic [width-1:0] d1,
     input logic [width-1:0] s,
     output logic [width-1:0] y)
    assign y = s ? d1 : d0;
endmodule
```

Test Bench: Overview & Concept (Simple w/ Asserts)

Adders

Half Adder

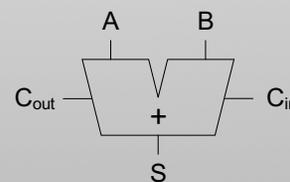


A	B	C_{out}	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

$$S = A \oplus B$$

$$C_{out} = AB$$

Full Adder

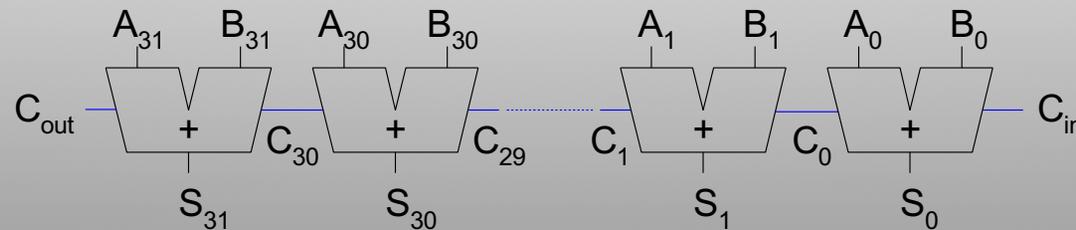


C_{in}	A	B	C_{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$S = A \oplus B \oplus C_{in}$$

$$C_{out} = AB + AC_{in} + BC_{in}$$

Ripple Carry Adder: Propagation Delay



Questions

Next Time

- Studio